

Recap: Chernoff-Hoeffding Bounds

$$Z = X_1 + \dots + X_n$$

$X_1, \dots, X_n$ mutually independent

$X_i \sim \text{Bernoulli}(p_i)$

$$\mu = \mathbb{E}[Z] = \sum_i \mathbb{E}[X_i] = \sum_i p_i$$

► (Chernoff/Hoeffding)

$$- \forall \delta > 0 \quad \mathbb{P}[Z \geq (1+\delta)\mu] \leq \left(\frac{e^\delta}{(1+\delta)^{1+\delta}} \right)^\mu \left. \vphantom{\mathbb{P}[Z \geq (1+\delta)\mu]} \right\} \leq e^{-\delta^2 \mu / 3}$$

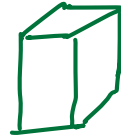
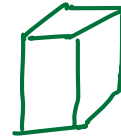
for $\delta \in (0, 1)$

$$- \forall \delta \in (0, 1) \quad \mathbb{P}[Z \leq (1-\delta)\mu] \leq \left(\frac{e^{-\delta}}{(1-\delta)^{1-\delta}} \right)^\mu$$

Application: Load Balancing



m jobs (balls)



n machines
(bins)

Assign jobs (randomly) so that
none of the machines have high load.

$E_j \equiv$ Machine j has load $\geq t$

Bound $P(\bigcup_j E_j)$

Ex: $P[A \cup B] \leq P[A] + P[B]$

Analyzing for $m = n$

$$X_{ij} = \begin{cases} 1 & \text{if job } i \text{ goes to machine } j \\ 0 & \text{o.w.} \end{cases} \quad \text{Bernoulli } (1/n)$$

$$Z_j = \# \text{ jobs on machine } j = \sum_i X_{ij}$$

$$\mathbb{E}[Z_j] = n \cdot \frac{1}{n} = 1$$

$$\mathbb{P}[Z_j \geq t] = \mathbb{P}[Z_j \geq (1 + (t-1)) \cdot 1] \leq \frac{e^{t-1}}{t^t}$$

$$\mathbb{P}[\exists j \ Z_j \geq t] \leq n \cdot \frac{1}{e} \cdot \left(\frac{e}{t}\right)^t$$

$$\leq n \cdot \frac{1}{e} \cdot \left(\frac{1}{n}\right)^2 \quad \text{for } t = \frac{3 \ln n}{\ln \ln n}$$

Power of two choices [Mitzenmacher]: $t = O(\ln \ln n)$

Random 3-SAT

Max 3-SAT : $\varphi \equiv C_1, C_2, \dots, C_m$ clauses in n vars

$$C_i = (\underbrace{l_{i_1}} \vee l_{i_2} \vee l_{i_3})$$

Can be a variable x_j or neg. $\bar{x}_j$.

e.g. $(x_1 \vee \bar{x}_9 \vee x_2), (x_5 \vee \bar{x}_2 \vee x_3), (\bar{x}_4 \vee \bar{x}_1 \vee \bar{x}_3)$

Goal : Assign $x_1 \dots x_n \in \{\text{True}, \text{False}\}$

To satisfy maximum no. of $C_1, \dots, C_m$

► For a random 3-SAT instance φ

(with $m \geq c \cdot n$ clauses for $c \geq 10/\epsilon^2$)

$$\mathbb{P}[\text{optimal assignment satisfies} \geq (\frac{7}{8} + \epsilon) \cdot m \text{ clauses}] \leq e^{-n}$$

Defining a random instance

Fix : $S_1, \dots, S_m \subseteq [n]$ $S_i = \{i_1, i_2, i_3\}$ (say)

$$l_{ij} = \begin{cases} x_{ij} & \text{w.p. } 1/2 \\ \bar{x}_{ij} & \text{w.p. } 1/2 \end{cases}$$

$$C_i = (l_{i_1} \vee l_{i_2} \vee l_{i_3})$$

Fix an assignment $\alpha \in \{0, 1\}^n \equiv \{\text{true}, \text{false}\}^n$

$$Z_\alpha = \sum_{i=1}^m Y_{i,\alpha} \quad \text{assignment } \alpha \text{ satisfies } i^{\text{th}} \text{ clause}$$

$$\mathbb{E}[Z_\alpha] = \sum_i \mathbb{E}[Y_{i,\alpha}] = m \cdot \frac{7}{8}$$

Union bounds

$$Z_\alpha = \sum_{i=1}^3 \mathbb{1}_{i,\alpha} \quad \text{assignment } \alpha \text{ satisfies } i^{\text{th}} \text{ clause}$$

$$\begin{aligned} \mathbb{P}[Z_\alpha \geq (\frac{7}{8} + \epsilon) m] &= \mathbb{P}\left[Z_\alpha \geq \left(1 + \epsilon \cdot \frac{8}{7}\right) \cdot \frac{7m}{8}\right] \\ &\leq e^{-\epsilon^2 \cdot \frac{8}{7} \cdot \frac{3}{8}} \leq e^{-2\epsilon^2 m} \quad \text{for } m \geq \frac{10}{\epsilon^2} \end{aligned}$$

$$\begin{aligned} \mathbb{P}[\exists \text{ assignment } \alpha \text{ s.t. } Z_\alpha \geq (\frac{7}{8} + \epsilon) m] \\ \leq 2^n \cdot e^{-2\epsilon^2 m} \leq e^{-n} \end{aligned}$$

Interlude: 0/1 random variables vs ± 1 variables

$$Z = Y_1 + \dots + Y_n$$

$$Y_i = \begin{cases} 1 & \text{w.p. } 1/2 \\ -1 & \text{w.p. } 1/2 \end{cases} \quad \text{i.i.d.}$$

Rademacher r.v.s

$$\mathbb{P}[|Z| \geq t] = ?$$

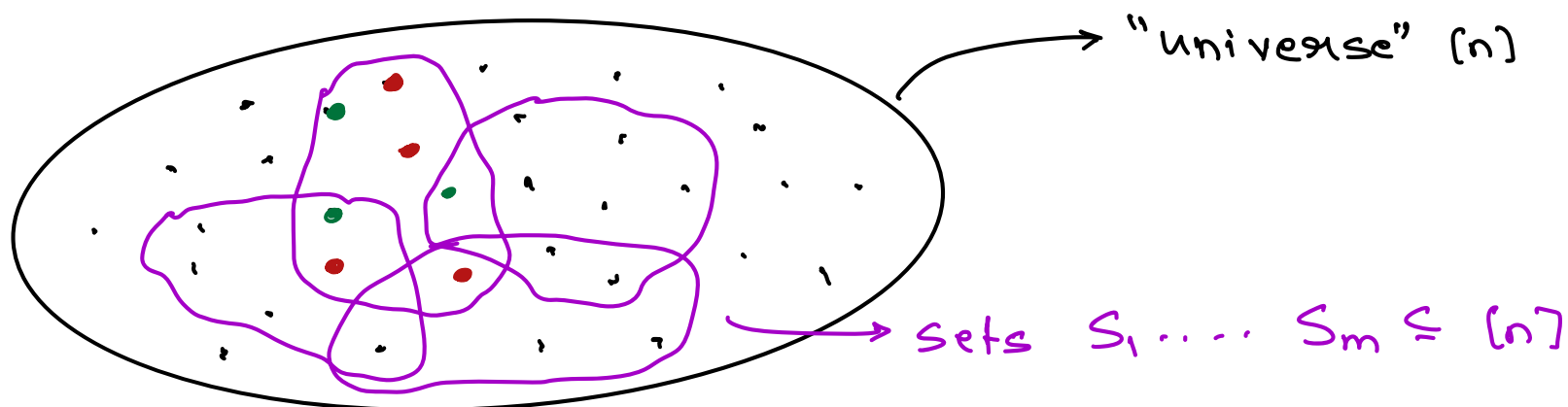
Consider $X_i = \frac{1 - Y_i}{2}$

$$X_i \sim \text{Bernoulli}(1/2)$$

Say $Z' = X_1 + \dots + X_n$

$$\begin{aligned} \mathbb{P}[|Z| \geq t] &= \mathbb{P}\left[\left|Z' - \frac{n}{2}\right| \geq \frac{t}{n}, \frac{n}{2}\right] \\ &\leq 2 \cdot e^{-(t^2/n^2) \cdot \frac{n}{2} \cdot \frac{1}{3}} = 2 \cdot e^{-t^2/6n} \end{aligned}$$

Balanced colorings (minimizing discrepancy)



Goal: Color points in $[n] = \{1, \dots, n\}$ by red/green
s.t. each S_i is as "balanced" as possible

$$x_j = \begin{cases} 1 & \text{if } j^{\text{th}} \text{ point is red} \\ -1 & \text{if } j^{\text{th}} \text{ point is green} \end{cases}$$

$$Z_i = \sum_{j \in S_i} x_j$$

$|Z_i|$ = "imbalance"

Balancing all sets

$$\mathbb{P}[|Z_i| \geq t] \leq 2 \cdot e^{-t^2/6 \cdot |S_i|} \leq 2 \cdot e^{-t^2/6n}$$

$$\mathbb{P}[\exists i \ |Z_i| \geq t] \leq n \cdot 2 \cdot e^{-t^2/6n}$$

$$\therefore \exists \text{ coloring with maximum imbalance } \leq \sqrt{12 n \ln n}$$

Alternate forms: Hoeffding's inequality

► For $x_1, \dots, x_n$ i.i.d. Rademacher variables

$$\mathbb{P}[|\sum a_i x_i| \geq t] \leq 2 \cdot e^{-t^2/2\|a\|^2}$$

Proof: $\mathbb{P}[\sum a_i x_i \geq t] = \mathbb{P}[e^{\lambda \sum a_i x_i} \geq e^{\lambda t}]$

$$\leq \frac{\mathbb{E} e^{\lambda \sum a_i x_i}}{e^{\lambda t}} = \frac{\prod_i \mathbb{E} e^{\lambda a_i x_i}}{e^{\lambda t}}$$

$$\mathbb{E} e^{\lambda a_i x_i} = \frac{1}{2} (e^{\lambda a_i} + e^{-\lambda a_i}) \leq e^{\lambda^2 a_i^2 / 2}$$

$$\mathbb{P}[\sum a_i x_i \geq t] \leq \frac{\prod_i e^{\lambda^2 a_i^2 / 2}}{e^{\lambda t}} = \frac{e^{\lambda^2 \|a\|^2 / 2}}{e^{\lambda t}}$$

$$\text{Optimal } \lambda = \frac{t}{\|a\|^2}$$

Hoeffding for bounded random variables

▷ Let $x_1, \dots, x_n$ be independent with $x_i \in [a_i, b_i]$

$$\mathbb{P} \left[\sum x_i - \mathbb{E}[\sum x_i] \geq t \right] \leq \exp \left(\frac{-2t^2}{\sum_i (b_i - a_i)^2} \right)$$

Ex: Prove above using $\mathbb{E} \left[e^{\lambda(x_i - \mathbb{E}[x_i])} \right] \leq \exp \left(\frac{\lambda^2 (b_i - a_i)^2}{8} \right)$